

# Analyse d'erreur de récurrences linéaires à l'aide de séries génératrices

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[arXiv:2011.00827](https://arxiv.org/abs/2011.00827) [math.NA]

# Polynômes de Legendre

$$(n + 1) P_{n+1}(x) = (2n + 1)x P_n(x) - n P_{n-1}(x)$$

$$x = \diamond(17/18)$$

	Intervalles	Flottants
n = 0	1.0000000000000000	1.0000000000000000
10	$[-0.386103889710 \pm 3.61e - 13]$	-0.386103889709935
20	$[ 0.29953942 \pm 5.12e - 9]$	0.299539415843941
30	$[-0.25184 \pm 7.04e - 6]$	-0.251843502730021
40	$[ 0.2 \pm 0.0280]$	0.214139934045693
50	$[\pm 56.2]$	-0.179433059907698
60	$[\pm 2.31e + 5]$	0.145569066400857
70	$[\pm 9.66e + 8]$	-0.112028696044130
80	$[\pm 4.09e + 12]$	0.0789787135424488
90	$[\pm 1.74e + 16]$	-0.0469035940021142
100	$[\pm 7.45e + 19]$	0.0164295146084927

# Un exemple jouet

[Boldo 2009]

$$c_{n+1} = 2c_n - c_{n-1} \quad (c_0 = \diamond(1/3), c_{-1} = 0)$$

	Intervalles	Flottants
n = 0	$[0.3333333333333333 \pm 1.49e - 17]$	0.3333333333333333
5	$[2.0000000000000000 \pm 3.78e - 15]$	2.0000000000000000
10	$[3.666666666666667 \pm 5.74e - 13]$	3.666666666666667
15	$[5.3333333333 \pm 5.29e - 11]$	5.333333333333334
20	$[7.00000000 \pm 1.60e - 9]$	7.000000000000001
25	$[8.666667 \pm 4.65e - 7]$	8.666666666666668
30	$[10.3333 \pm 4.41e - 5]$	10.33333333333333
35	$[12.000 \pm 8.82e - 4]$	12.00000000000000
40	$[1.4e + 1 \pm 0.406]$	13.666666666666667
45	$[\pm 21.3]$	15.333333333333334
50	$[\pm 5.04e + 2]$	17.00000000000000

# Analyse d'erreur naïve

(erreur absolue / “virgule fixe” pour simplifier)

$$c_{n+1} = 2c_n - c_{n-1}$$

$$\tilde{c}_{n+1} = \diamond(2\tilde{c}_n - \tilde{c}_{n-1})$$

$$= 2\tilde{c}_n - \tilde{c}_{n-1} + \varepsilon_n \quad |\varepsilon_n| \leq \mathbf{u}$$

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► récurrence :  $|\tilde{c}_n - c_n| \leq 3^n \mathbf{u}$

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► un poil plus fin :

$$|\tilde{c}_n - c_n| \leq (\lambda_+ \alpha_+^n + \lambda_- \alpha_-^n - 4) \mathbf{u} \approx 2.4^n \mathbf{u}$$

$$\alpha_{\pm} = 1 \pm \sqrt{2}, \quad \lambda_{\pm} = 4 \pm 3\sqrt{2}$$



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C'est à peu près ce que fait l'arithmétique d'intervalles.

# Analyse d'erreur plus raisonnable

("virgule fixe")

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$$(\delta_0 = \delta_1 = 0)$$

*erreur globale* ↗

↖ *erreur locale*

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$$|\delta_n| \leq \frac{n(n-1)}{2} \mathbf{u}$$

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Il faut transporter une expression linéaire en les  $\varepsilon_k$ .

Les calculs peuvent devenir indigestes.

# Un mot d'histoire

Peu de littérature sur les récurrences dans le sens direct

in the backward direction. There has been less attention devoted to computation which utilizes the difference equation in the forward direction, not because a forward algorithm is more difficult to analyze, but rather for the opposite reason—that its analysis was considered straightforward. Of the above

[Wimp 1972]

**von Neumann & Goldstine 1947, Turing 1948** — Systèmes triangulaires

**Henrici 1962** — Schémas numériques pour les EDO

**Oliver 1965** — Propagation des erreurs relatives dans les réc. linéaires

**Liu & Kaneko 1969** — Filtres linéaires en flottant (→ fns de transfert)

**Hilaire & Lopez 2013** — Filtres linéaires, pire cas

# Séries génératrices



$$(\mathbf{a}_n)_{n \in \mathbb{Z}} \longleftrightarrow \mathbf{a}(z) = \sum_{n=-\infty}^{\infty} \mathbf{a}_n z^n$$

$$\delta_{n+1} = 2 \delta_n - \delta_{n-1} + \varepsilon_n$$

$$\downarrow \sum_n \square z^n$$

$$z^{-1} \delta(z) = 2 \delta(z) - z \delta(z) + \varepsilon(z)$$

$$\delta(z) = \frac{z}{(1-z)^2} \varepsilon(z)$$

$$z \sum_n \mathbf{a}_n z^n = \sum_n \mathbf{a}_{n-1} z^n$$

- ▶ Formules (produit, composée, ...)
- ▶ Méthodes analytiques
  - ▶ Extraction de coefficients
  - ▶ Asymptotique
- ▶ Séries majorantes
- ▶ Calcul efficace
- ▶ ...

# Séries majorantes



On note  $f \ll \hat{f}$  lorsque  $\forall n, |f_n| \leq \hat{f}_n$ .  
( $f(z) = \sum_n f_n z^n, \hat{f}(z) = \sum_n \hat{f}_n z^n$ )

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$$\triangleright \text{si } f(z) \ll \hat{f}(z) \text{ et } g(z) \ll \hat{g}(z), \\ f(z) g(z) \ll \hat{f}(z) \hat{g}(z)$$

$$\left| [z^n] (f(z) g(z)) \right| = \left| \sum_{i+j=n} f_i g_j \right| \leq \sum_{i+j=n} \hat{f}_i \hat{g}_j = [z^n] (\hat{f}(z) \hat{g}(z))$$

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## Flottants, erreur relative

$$c_{n+1} = 2c_n - c_{n-1}$$

$$\tilde{c}_{n+1} = \diamond(2\tilde{c}_n - \tilde{c}_{n-1})$$

$$= (2\tilde{c}_n - \tilde{c}_{n-1})(1 + \epsilon_n) \quad |\epsilon_n| \leq \mathbf{u}$$

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$$(1 + \varepsilon_n) \quad \mathbf{c}_{n+1} = 2 \mathbf{c}_n - \mathbf{c}_{n-1}$$

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$$= (2 \tilde{\mathbf{c}}_n - \tilde{\mathbf{c}}_{n-1}) (1 + \varepsilon_n) \quad |\varepsilon_n| \leq \mathbf{u}$$

$$\delta_{n+1} - \mathbf{c}_{n+1} \varepsilon_n = (2 \delta_n - \delta_{n-1}) (1 + \varepsilon_n)$$

$$\delta_n = \tilde{\mathbf{c}}_n - \mathbf{c}_n$$

$$\delta_{n+1} - 2 \delta_n + \delta_{n-1} = \varepsilon_n (\mathbf{c}_{n+1} + 2 \delta_n - \delta_{n-1})$$

$$(z^{-1} - 2 + z) \delta(z) = \varepsilon(z) \odot (z^{-1} \mathbf{c}(z) + (2 - z) \delta(z))$$

$$\delta(z) = \frac{(z \varepsilon(z)) \odot (\mathbf{c}(z) + z(2 - z) \delta(z))}{(1 - z)^2}$$

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$$\#f(z) = \sum_n |f_n| z^n$$

$$c(z) \ll \hat{c}(z)$$

$$\delta(z) \ll \frac{z(2+z)\mathbf{u}}{(1-z)^2} \#f(z) + \frac{\hat{c}(z)\mathbf{u}}{(1-z)^2}$$

# Équations majorantes

## Lemme

[Cauchy ?]

Soient  $\hat{a}(z), \hat{b}(z) \in \mathbb{R}_+[[z]]$  avec  $\hat{a}(0) = 0$ . Supposons  $y \in \mathbb{R}_+[[z]]$  t.q.

$$y(z) \ll \hat{a}(z) y(z) + \hat{b}(z).$$

Alors  $y(z)$  est bornée par la solution de  $\hat{y}(z) = \hat{a}(z) \hat{y}(z) + \hat{b}(z)$ , i.e.,

$$y(z) \ll \hat{y}(z) = \frac{\hat{b}(z)}{1 - \hat{a}(z)}.$$

**Preuve.** ▶  $y_0 \leq \hat{b}_0 = \hat{y}_0$

$$\text{▶ } y_n \leq \sum_{i=0}^n \hat{a}_i y_{n-i} + \hat{b}_n$$

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$$\text{▶ } y_n \leq \sum_{i=1}^n \hat{a}_i y_{n-i} + \hat{b}_n \leq \sum_{i=1}^n \hat{a}_i \hat{y}_{n-i} + \hat{b}_n = \hat{y}_n$$

## Borne sur l'erreur en flottants

$$\# \delta(z) \ll \underbrace{\frac{z(2+z)\mathbf{u}}{(1-z)^2}}_{\hat{\mathbf{a}}(z)} \# \delta(z) + \underbrace{\frac{\hat{\mathbf{c}}(z)\mathbf{u}}{(1-z)^2}}_{\hat{\mathbf{b}}(z)}$$

$$\begin{aligned} \Delta(z) &\ll \frac{\hat{\mathbf{b}}(z)}{1 - \hat{\mathbf{a}}(z)} \\ &= \frac{\hat{\mathbf{c}}(z)\mathbf{u}}{1 - 2(1+\mathbf{u})z + (1-\mathbf{u})z^2} \\ &= \frac{\hat{\mathbf{c}}(z)\mathbf{u}}{(1-\alpha z)(1-\beta z)} \\ &\ll \frac{|c_0|\mathbf{u}}{(1-\alpha z)^4} \end{aligned}$$

→

$$\begin{aligned} c(z) &= \frac{c_0}{(1-z)^2} \\ &\ll \frac{|c_0|}{(1-z)^2} \end{aligned}$$

Exponentiel, mais  $O(1)$   
si  $n = O(\mathbf{u}^{-1/2})$

$$\alpha = 1 + \sqrt{3\mathbf{u}} + O(\mathbf{u})$$

$$\alpha \leq 1 + 2\sqrt{\mathbf{u}} \quad \text{pour } \mathbf{u} \leq 0.008$$

$$|\delta_n| \leq \frac{|c_0|}{6} (n+3)^3 \alpha^n \mathbf{u}$$

( $\delta_n =$  erreur absolue sur  $c_n$ )

## D'autres exemples

### Polynômes de Legendre

▶  $P_{n+1}(x) = \frac{1}{n+1} [(2n+1)xP_n(x) - nP_{n-1}(x)]$

▶ En arithmétique virgule fixe :

$$|\tilde{p}_n - P_n(x)| \leq \frac{3}{4} (n+1)(n+2) u \quad (-1 \leq x \leq 1)$$

▶ Coefficients polynomiaux  $\rightarrow$  équ. diff.

Utile pour le calcul de règles de quadrature de Gauss-Legendre [Johansson-M. 2018]

### Corde vibrante [Boldo 2009, Boldo *et al.* 2013, 2014]

▶  $\frac{\partial^2 p}{\partial t^2} - c^2 \frac{\partial^2 p}{\partial x^2} = 0 \quad \xrightarrow{\text{discr.}} \quad p_i^{k+1} = 2p_i^k - p_i^{k-1} + \alpha (p_{i+1}^k - 2p_i^k + p_{i-1}^k)$

▶ En double :

$$|\tilde{p}_i^k - p_i^k| \leq 39 \cdot 2^{-52} \cdot (k+1)(k+2)$$

▶ Normes sur  $\Omega = \mathbb{R}[x] / \langle x^{2n} - 1 \rangle$  et séries dans  $\Omega[[t]]$

# Nombres de Bernoulli

$$B_n = 1, \frac{-1}{2}, \frac{1}{6}, 0, \frac{-1}{30}, 0, \frac{1}{42}, 0, \frac{-1}{30}, 0, \frac{5}{66}, 0, \frac{-691}{2730}, 0, \frac{7}{6}, 0, \frac{-3617}{510}, \dots \quad |B_{2k}| \sim \frac{2(2k)!}{(2\pi)^{2k}}$$

$$b_k = \frac{B_{2k}}{(2k)!} \quad b(z) = \sum_{k=0}^{\infty} b_k z^k = \frac{\sqrt{z}/2}{\tanh(\sqrt{z}/2)}$$

## Algorithmme

[Brent 1980, sugg. de Reinsch]

$$b_k = \frac{1}{(2k)! 4^k} - \sum_{j=0}^{k-1} \frac{b_j}{(2k+1-2j)! 4^{k-j}}$$

be used with sufficient guard digits, or **a more stable recurrence** must be used. If we multiply both sides of (30) by  $\sinh(x/2)/x$  and equate coefficients, we get the recurrence

$$C_k + \frac{C_{k-1}}{3! 4} + \dots + \frac{C_1}{(2k-1)! 4^{k-1}} = \frac{2k}{(2k+1)! 4^k} \quad (36)$$

If (36) is used to evaluate  $C_k$ , using precision  $n$  arithmetic, **the error is only  $O(k^2 2^{-n})$** . Thus,

[Brent 1980]

## Erreur relative dans le calcul de $b_k$

**Exercice 4.35** Prove (or give a plausibility argument for) the statements made in §4.7 that: (a) if a recurrence based on (4.59) is used to evaluate the scaled Bernoulli number  $C_k$ , using precision  $n$  arithmetic, then the relative error is of order  $4^k 2^{-n}$ ; and (b) if a recurrence based on (4.60) is used, then the relative error is  $O(k^2 2^{-n})$ .

[Brent & Zimmermann 2010]

### Conjecture

[Brent & Zimmermann ?]

On a  $\tilde{b}_k = b_k (1 + \eta_k)$  où  $\eta_k = O(k \cdot u)$ .

$\tilde{b}_k$  = valeur calculée en flottants,  $u$  = unité d'arrondi

### Remarque

À comprendre comme  $\eta_k = O(k \cdot u)$  pour  $k = O(u^{-1})$ ,  
ou encore  $|\eta_k| \leq C_k u$  quand  $u \rightarrow 0$  où  $C_k = O(k)$  (resp.  $O(k^2)$ )



## Borne explicite sur l'erreur relative

$$b(z) = \frac{\sqrt{z}/2}{\tanh(\sqrt{z}/2)} \quad b_k = \frac{1}{(2k)! 4^k} - \sum_{j=0}^{k-1} \frac{b_j}{(2k+1-2j)! 4^{k-j}}$$

$$\tilde{b}_k = b_k (1 + \eta_k)$$

### Proposition

L'erreur relative sur le résultat satisfait

$$|\eta_k| \leq (1 + 21.2 u)^k (1.1 k + 446) u$$

Corollaire :  $|\eta_k| \leq (3k + 1213) u$  pour  $u < 2^{-16}$  et  $43ku \leq 1$ .

# Schéma de la preuve

## Analyse d'erreur locale.

$$\tilde{b}_k = \frac{1 + \mathbf{s}_k}{(2k)! 4^k} - \sum_{j=0}^{k-1} \frac{\tilde{b}_j (1 + \mathbf{t}_{k,j})}{(2k+1-2j)! 4^{k-j}}$$

$\tilde{b}_k =$  valeurs calculées

$$\begin{aligned} |\mathbf{s}_k| &\leq \hat{\theta}_{2k} \\ |\mathbf{t}_{k,j}| &\leq \hat{\theta}_{3(k-j)+2} \\ \text{où } \hat{\theta}_n &= (1 + u)^n - 1 \end{aligned}$$

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**Linéarité.**  $\delta_k := \tilde{b}_k - b_k =$  erreur globale

$$\delta_k = \frac{\mathbf{s}_k}{(2k)! 4^k} - \sum_{j=0}^{k-1} \frac{\delta_j + \tilde{b}_j \mathbf{t}_{k,j}}{(2k+1-2j)! 4^{k-j}}$$

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## Inéquation sur l'erreur globale.

$$\delta(z) \ll \check{S}(z) \tilde{C}(z) + \check{S}(z) \tilde{S}(z) \# b(z) + \check{S}(z) \tilde{S}(z) \# \delta(z) \quad \# f(z) = \sum_k |f_k| z^k$$

$$\begin{aligned} \text{où } C(z) &= \cosh(\sqrt{z}/2), & S(z) &= (\sqrt{z}/2)^{-1} \sinh(\sqrt{z}/2), & \check{S}(z) &= \frac{(\sqrt{z}/2)}{\sin(\sqrt{z}/2)}, \\ \tilde{C}(z) &= C(\alpha^2 z) - C(z), & \tilde{S}(z) &= S(\alpha^4 z) - S(z) - (\alpha^2 - 1) & \text{avec } \alpha &= 1+u \end{aligned}$$

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**Majorant « explicite ».**

$$\delta(z) \ll \frac{\check{S}(z) \tilde{C}(z) + \check{S}(z) \tilde{S}(z) \#b(z)}{1 - \check{S}(z) \tilde{S}(z)} =: \hat{\delta}(z)$$

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**Asymptotique.**

$$\hat{\delta}(z) = \left( \frac{2(1 - \cosh w) \cos(w)}{w^{-2} \sin(w)^2} + \frac{4(\cosh w - 1) + w \sinh w}{w^{-1} \sin w} \right) u + O(u^2)$$

$w = \sqrt{z}/2$

Unique pôle de module minimal en  $z = 4\pi^2$ ,  
multiplicité (en  $z$ ) = 2

$$\Rightarrow \hat{\delta}_k = O(k(2\pi)^{-2k}) \cdot u + O(u^2)$$

$$\Rightarrow \eta_k = "O(k \cdot u)"$$

# Schéma de la preuve

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## Contrôle du pôle dominant.

Supposons  $u \leq 2^{-16}$ .

Alors  $\hat{\delta}(z)$  a un pôle en  $\gamma = \left( \frac{2\pi}{1 + \varphi(u)} \right)^2$  où  $0 \leq \varphi(u) \leq 2(\cosh \pi - 1)u$ .

C'est le seul avec  $|z| < 153.7 \approx (3.9\pi)^2$ .

(Un peu d'analyse, et comparaison avec le cas limité par le théorème de Rouché.)

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## Symbolique-numérique.

$$\hat{\delta}(z) \ll \frac{2|\mathbf{R}(u) - 1|}{1 - z/\gamma} + \frac{2\varphi(u)(2 + \varphi(u))z/(2\alpha)^2}{(1 - z/\gamma)^2} + \frac{\mathbf{A}_\lambda(u)}{1 - z/(\lambda\gamma)}$$

Formule de Cauchy + arithmétique d'intervalles.



# Sommes partielles de séries D-finies

Contexte : résolution numérique de  $L \cdot u = 0$  où  $L \in \mathbb{C}[z]\langle d/dz \rangle$



Calculer un encadrement de  $\sum_{n=0}^{N-1} u_n \zeta^n$ .

$L, u_{0:r-1}, \zeta, N$  donnés

► **Précision cible**  $p$

► **Hypothèses**

point ordinaire —  $a_r(0) \neq 0$   
cvgce « évidente » —  $|\zeta| < \min \{|\xi| : a_r(\xi) = 0\}$   
cvgce géométrique —  $N = \Theta(p)$

$L = a_r(z) \left(\frac{d}{dz}\right)^r + \dots$

# Naïve approach

Recurrence step:

(naïve interval arithmetic, with  $u_n \approx \Theta(1)$ )

$$u_n = \frac{-1}{b_s(n)} [b_{s-1}(n) u_{n-1} + \dots + b_1(n) u_{n-s+1} + b_0(n) u_{n-s}]$$

$\text{rad} \approx \rho$

$$\text{rad} \gtrsim \left( \sum_i \left| \frac{b_i(s)}{b_s(n)} \right| \right) \rho \gtrsim s \rho$$

$$\text{rad}(u_n) = 2^{\Theta(n)}$$

$$\text{rad}(\sum^N u_n \zeta^n) = 2^{\Theta(N)} \text{ (unless } \zeta \text{ small)}$$

Accuracy target  $2^{-p} \Rightarrow$  Need  $\Omega(N) = \Omega(p)$  guard bits



(This is not a numerical stability issue.)

# Running Error Analysis

$$\mathbf{u}_n = \frac{-1}{b_s(n)} [b_{s-1}(n) \mathbf{u}_{n-1} + \cdots + b_1(n) \mathbf{u}_{n-s+1} + b_0(n) \mathbf{u}_{n-s}]$$

$$\tilde{\mathbf{u}}_n = \frac{-1}{b_s(n)} [b_{s-1}(n) \tilde{\mathbf{u}}_{n-1} + \cdots + b_1(n) \tilde{\mathbf{u}}_{n-s+1} + b_0(n) \tilde{\mathbf{u}}_{n-s}] + \varepsilon_n$$

$\tilde{\mathbf{u}}_n$  = computed sequence (e.g. floating-point)

The global error  $\delta_n = \tilde{\mathbf{u}}_n - \mathbf{u}_n$  satisfies

$$b_s(n) \delta_n + b_{s-1}(n) \delta_{n-1} + \cdots + b_0(n) \delta_{n-s} = b_s(n) \varepsilon_n$$

local error,  
**known** bound  $|\varepsilon_n| \leq \hat{\varepsilon}_n$   
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Therefore:

$$a_r(z) \delta^{(r)}(z) + \dots + a_1(z) \delta'(z) + a_0(z) \delta(z) = Q(z \, d/dz) \varepsilon(z)$$

$$\delta(z) = \sum_n \delta_n z^n, \quad \varepsilon(z) = \sum_n \varepsilon_n z^n$$

Compute a **bound** on  $\delta_n$  given one on  $\varepsilon_n$ ?

# Global Error

$$\alpha_r(z) \delta^{(r)}(z) + \dots + \alpha_0(z) \delta(z) = Q(z \frac{d}{dz}) \varepsilon(z) \xrightarrow{\text{maj.}} \hat{\delta}'(z) = \hat{a}(z) \hat{\delta}(z) + \hat{\varepsilon}(z)$$

$$|\delta_0|, \dots, |\delta_{r-1}| \leq \hat{\delta}_0 \Rightarrow \delta(z) \ll \hat{\delta}(z)$$

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$$|\delta_0|, \dots, |\delta_{r-1}| \leq \hat{\delta}_0 \Rightarrow \delta(z) \ll \hat{\delta}(z)$$

$$\hat{\delta}(z) = \hat{h}(z) \left( \text{cst} + \int_0^z \frac{\hat{\varepsilon}(w)}{\hat{h}(w)} dw \right), \quad \hat{h}(z) = \exp \int_0^z \hat{a}(z) dz$$

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take  $\hat{\varepsilon}_n = \bar{\varepsilon} \hat{h}_n$       $\bar{\varepsilon} \lesssim n^r \mathbf{u}$  since  $|u_n| \lesssim \hat{h}_n$



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$$|\tilde{\mathbf{u}}(\zeta) - \mathbf{u}(\zeta)| \leq \hat{\delta}(|\zeta|) = O(\bar{\varepsilon}) = O(2^{-p_{\text{work}} + \log N})$$

#guard bits =  $O(\log N)$  for fixed  $L, \text{ini}, \zeta$

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$$\# \text{guard bits} = O(\log N) \quad \text{for fixed } L, \text{ini}, \zeta$$



The same computation yields a bound on the **truncation** error!

(Replace  $\varepsilon(z)$  by a residual accounting for the neglected tail.)

# Practical Issues



- The Cauchy majorants are *far* too coarse
  - ▶ Use sharper variants [M. 2019]
- Simple majorants cannot be sharp for small  $n$ 
  - ▶ Switch from interval summation to running error analysis
- Need to choose the working precision ( $\leftrightarrow$  cutoff point) in advance
  - ▶ Heuristics based on asymptotics...
- Good choice of  $\hat{\varepsilon}(z)$  (tight & easily computable) not clear

Current status: works well for (some) large equations met in practice, but sometimes slower than naïve interval summation.

# Conclusion



Une analyse d'erreur pénible ?  
Essayez les séries génératrices !

- ▶ Erreur locale  $\leftrightarrow$  erreur globale
- ▶ Expressions exactes ou équations
- ▶ Séries majorantes



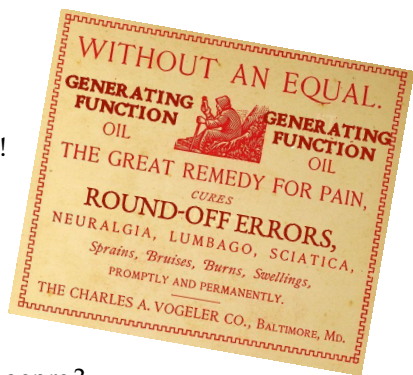
Qui a déjà vu des choses de ce genre ?



Algorithme pratique pour les D-finies  
(y compris cas singulier)



Réurrences "à rebours" (Miller) ?  
Polynômes orthogonaux généraux ?  
Schémas d'intégration numérique ?



# Polynômes de Legendre

[Johansson & M. 2018]

$$p_{n+1} = \frac{1}{n+1} [(2n+1)x p_n - n p_{n-1}]$$

$x$  fixé,  $p_n = P_n(x)$

$$\tilde{p}_{n+1} = \frac{1}{n+1} [(2n+1)x \tilde{p}_n - n \tilde{p}_{n-1}] + \varepsilon_{n+1}$$

$|\varepsilon_n| \leq 3u$  (FxP)

$$\delta_n = \tilde{p}_n - p_n$$

$$(n+1) \delta_{n+1} = (2n+1)x \delta_n - n \delta_{n-1} + (n+1) \varepsilon_{n+1}$$

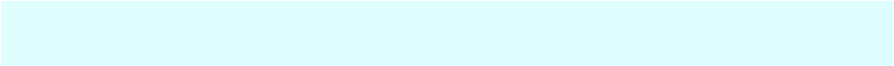
$$(1 - 2xz + z^2) \delta'(z) = z(x - z) \delta(z) + \varepsilon'(z)$$

$$\delta(z) = \frac{1}{\sqrt{1 - 2xz + z^2}} \int_0^z \frac{\varepsilon'(z)}{\sqrt{1 - 2xz + z^2}}$$

$$z \sum_n a_n z^n = \sum_n a_{n-1} z^n$$

$$\frac{d}{dz} \left( \sum_n f_n z^n \right) = \sum_n (n+1) f_{n+1} z^n$$

## Majoration de la solution

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$$\frac{3u}{(1-z)^2}$$

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►  $\frac{1}{\sqrt{1-2xz+z^2}} = \frac{1}{\sqrt{(1-e^{i\theta}z)(1-e^{-i\theta}z)}}$   $x = \cos \theta$

► or  $\frac{1}{\sqrt{1-e^{\pm i\theta}z}} \ll \frac{1}{\sqrt{1-z}}$  par composition de  $\frac{1}{\sqrt{1-z}} \in \mathbb{R}_+[[z]]$  et  $e^{i\theta}z$

► donc  $\frac{1}{\sqrt{1-2xz+z^2}} \ll \frac{1}{1-z}$

(En fait  $\frac{1}{\sqrt{1-2xz+z^2}} = \sum_n P_n(x) z^n$ , et il est classique que  $|P_n(x)| \leq 1$ .)



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(En fait  $\frac{1}{\sqrt{1-2xz+z^2}} = \sum_n P_n(x) z^n$ , et il est classique que  $|P_n(x)| \leq 1$ .)

## Majoration de la solution

$$\delta(z) = \frac{1}{\sqrt{1-2xz+z^2}} \int_0^z \frac{\varepsilon'(z)}{\sqrt{1-2xz+z^2}} \quad \varepsilon(z) \ll \frac{3u}{1-z}$$

$$\delta(z) \ll \frac{1}{1-z} \int \frac{3u}{(1-z)^2} \frac{1}{1-z} = \frac{3}{2} \frac{1}{(1-z)^3} u$$

►  $\frac{1}{\sqrt{1-2xz+z^2}} = \frac{1}{\sqrt{(1-e^{i\theta}z)(1-e^{-i\theta}z)}} \quad x = \cos \theta$

► or  $\frac{1}{\sqrt{1-e^{\pm i\theta}z}} \ll \frac{1}{\sqrt{1-z}}$  par composition de  $\frac{1}{\sqrt{1-z}} \in \mathbb{R}_+[[z]]$  et  $e^{i\theta}z$

► donc  $\frac{1}{\sqrt{1-2xz+z^2}} \ll \frac{1}{1-z}$

(En fait  $\frac{1}{\sqrt{1-2xz+z^2}} = \sum_n P_n(x) z^n$ , et il est classique que  $|P_n(x)| \leq 1$ .)

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