

# The Dynamic Dictionary of Mathematical Functions

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Numerical Analysis Seminar, Oxford, November 20th, 2014

# **1** The DDMMF

# The DDMF Project

(~2009–2012)

**D.** ddmf.msr-inria.inria.fr/1.9.1/ddmf

## Dynamic Dictionary of Mathematical Functions

Welcome to this interactive site on [Mathematical Functions](#), with properties, truncated expansions, numerical evaluations, plots, and more. The functions currently presented are elementary functions and special functions of a single variable. More functions — special functions with parameters, orthogonal polynomials, sequences — will be added with the project advances.

**This is release 1.9.1 of DDMF**  
[Select a mathematical rendering](#) to enable access to the contents

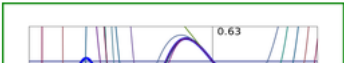
**What's new?** The main changes in this release 1.9.1, dated May 2013, are:

- Proofs related to Taylor polynomial approximations.

Release [history](#).

**More on the project:**

- [Help](#) on selecting and configuring the mathematical rendering
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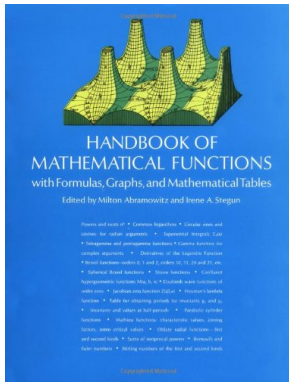


### Mathematical Functions

- The Airy function of the first kind
- The Airy function of the second kind
- The Anger function
- The inverse cosine
- The inverse hyperbolic cosine
- The inverse cotangent
- The inverse hyperbolic cotangent
- The inverse cosecant
- The inverse hyperbolic cosecant
- The inverse secant
- The inverse hyperbolic secant
- The inverse sine
- The inverse hyperbolic sine
- The inverse tangent
- The inverse hyperbolic tangent
- The modified Bessel function of the first kind
- The modified Bessel function of the second kind
- The Bessel function of the first kind
- The Chebyshev function of the first kind
- The Chebyshev function of the second kind
- The hyperbolic cosine integral
- The cosine integral
- The cosine
- The hyperbolic cosine
- The Coulomb function
- The Whittaker's parabolic function

**Joint work with:**  
Alexandre Benoit  
**Frédéric Chyzak**  
Alexis Darrasse  
Stefan Gerhold  
Thomas Grégoire  
Stéphane Henriot  
Christoph Koutschan  
Sébastien Maulat  
Bruno Salvy  
Shiv Shankar

# An Encyclopedia of Special Functions



## Ascending Series

$$10.4.2 \quad \text{Ai}(z) = c_1 f(z) - c_2 g(z)$$

$$10.4.3 \quad \text{Bi}(z) = \sqrt{3} [c_1 f(z) + c_2 g(z)]$$

$$f(z) = 1 + \frac{1}{3!} z^3 + \frac{1 \cdot 4}{6!} z^6 + \frac{1 \cdot 4 \cdot 7}{9!} z^9 + \dots$$

$$= \sum_0^{\infty} 3^k \left(\frac{1}{3}\right)_k \frac{z^{3k}}{(3k)!}$$

$$g(z) = z + \frac{2}{4!} z^4 + \frac{2 \cdot 5}{7!} z^7 + \frac{2 \cdot 5 \cdot 8}{10!} z^{10} + \dots$$

$$= \sum_0^{\infty} 3^k \left(\frac{2}{3}\right)_k \frac{z^{3k+1}}{(3k+1)!}$$

$$\left(\alpha + \frac{1}{3}\right)_0 = 1$$

$$3^k \left(\alpha + \frac{1}{3}\right)_k = (3\alpha+1)(3\alpha+4) \dots (3\alpha+3k-2)$$

( $\alpha$  arbitrary;  $k=1, 2, 3, \dots$ )

(See 6.1.22.)

## 10.4.4

$$c_1 = \text{Ai}(0) = \text{Bi}(0) / \sqrt{3} = 3^{-2/3} / \Gamma(2/3)$$

$$= .35502 \ 80538 \ 87817$$

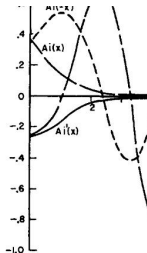


FIGURE 10





# An Encyclopedia of Special Functions

```
Untitled (1)* - [Server 1] - Maple 17
File Edit View Insert Format Table Drawing Plot Spreadsheet Tools
> evalf[30](BesselJ(3, 1));
0.0195633539826684059189053216218
> int(AiryAi(t)*exp(-t^3/x), t=0..infinity)
assuming x > 0;
1/9 e^(1/18 x) (x BesselK(11/3, 1/3) exp(x/18) + 5/3)
```

7.2(1) Error Functions

7.2.1 
$$\operatorname{erf} z = \frac{2}{\sqrt{\pi}} \int_0^z e^{-t^2} dt,$$

7.2.2 
$$\operatorname{erfc} z = \frac{2}{\sqrt{\pi}} \int_z^\infty e^{-t^2} dt = 1 - \operatorname{erf} z,$$

# Demo!

D. ddmf.msr-inria.inria.fr/1.9.1/ddmf

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- The Whittaker's parabolic function

# Behind the Scenes...

- **DynaMoW** — Ocaml + Maple + web [Chyzak & Darrasse]
- **AlgoLib**
  -  <http://algo.inria.fr/libraries/>  
GNU LGPL
  - **gfun** — D-finite power series [Salvy & Zimmermann]
  - **MGfun** — Multivariate D-finite functions [Chyzak]
  - **NumGfun** — D-finite analytic functions [M.]
  - ...
- (+ DDMF-specific Maple libraries)



## **2** Identities

# D-Finite Functions

[Stanley 1980, ...]

An analytic function  $y: \mathbb{C} \rightarrow \mathbb{C}$  is **D-finite** (holonomic) iff it satisfies a linear homogeneous ODE with polynomial coefficients:

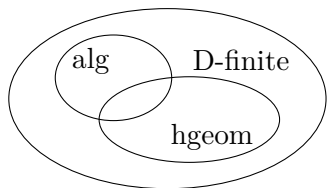
$$a_r(z) y^{(r)}(z) + \cdots + a_1(z) y'(z) + a_0(z) y(z) = 0, \quad a_j \in \mathbb{K}[z]$$

**Examples:**  $\sin$  “=”  $\{y''(z) + y(z) = 0, \quad y(0) = 0, \quad y'(0) = 1\}$

$K_0$  “=”  $\{z y''(z) + y'(z) - z y(z) = 0, \quad (\text{ini?})\}$

**More examples:**  $\exp$ ,  $\ln$ ,  $Ai$ ,  $\text{erf}$ ,  $J_\nu$ ,  
algebraic fns, hypergeometric fns...

**Not D-finite:**  $\zeta$ ,  $\Gamma$ ,  $W$ ,  $\tan$ ,  $C_{a,q}$ , ...



# D-Finite Functions

[Stanley 1980, ...]

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$$a_r(z) y^{(r)}(z) + \cdots + a_1(z) y'(z) + a_0(z) y(z) = 0, \quad a_j \in \mathbb{K}[z]$$

**Equivalently:**

$$\dim_{\mathbb{K}(z)} \text{Span}_{\mathbb{K}(z)} (D^i \cdot y)_{i \in \mathbb{N}} < \infty$$

A light blue rectangular box containing a diagram of the span of derivatives of a function y. The diagram shows a vertical list of terms:

- $D \cdot y = y$  (with y centered above the equals sign)
- $D^2 \cdot y = y'$  (with y' centered above the equals sign)
- $D^3 \cdot y = -\frac{a_0}{a_3} y - \frac{a_1}{a_3} y' - \frac{a_2}{a_1} y''$  (with y, y', and y'' centered above their respective terms)
- $D^4 \cdot y = *y + *y' + *y''$  (with y, y', and y'' centered above their respective terms)
- $\vdots$

# Closure Properties

D-finite functions over  $\mathbb{K}$  form a  $\mathbb{K}$ -algebra.

Example:

$$f(z) = \sin z$$

$$f''(z) + f(z) = 0$$

$$g(z) = \text{Ai}(z)$$

$$g''(z) - zg(z) = 0$$

$$fg = fg$$

$$(fg)' = f'g + fg'$$

$$\begin{aligned}(fg)'' &= f''g + f'g' + f'g' + fg'' \\ &= -fg + 2f'g' + zfg \\ &= (z-1)fg + 2f'g'\end{aligned}$$

$$(fg)^{(3)} = fg + (3z-1)f'g + (z-3)fg'$$

$$\begin{aligned}(fg)^{(4)} &= (z^2 - 6z + 1)fg + 4f'g \\ &\quad + 2fg' + (4z - 4)f'g'\end{aligned}$$

5 vectors

dimension 4

$$(z+1)(fg)^{(4)} - (fg)^{(3)} - 2(z^2+1)(fg)'' - (z+5)(fg)' + (z^3+\dots)(fg) = 0$$

# Proof of Identities

**Proposition.**  $\cos^2 z + \sin^2 z = 1$ .

**Proof.**  $\cos^2 z + \sin^2 z - 1 = O(z^9)$ . □

Why is this a proof?

- $\cos$  and  $\sin$  both satisfy equations of order 2
- hence,  $\cos^2$  and  $\sin^2$  satisfy equations of order  $\leq 4$
- 1 satisfies  $y' = 0$
- hence,  $\cos^2 + \sin^2 - 1$  satisfies an equation of order  $\leq 9$

(Actually,  $\cos^2 x$ ,  $\sin^2 x$  and 1 all satisfy  $y^{(3)} + 4y' = 0$ .)

In CS terms: equality is decidable.

# Generalisations

[Lipshitz, Zeilberger, Chyzak & Salvy, ...]

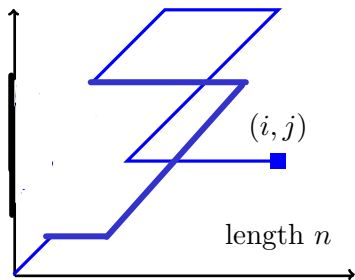
$$T_n(x) \quad (-1)^{m-k} k! \binom{n-k}{m-k} \left\{ \begin{matrix} n+1 \\ k+1 \end{matrix} \right\} \quad \frac{\cos(zt)}{\sqrt{1-t^2}}$$

$$\binom{n}{k} \quad \frac{\exp\left(\frac{4u(xy - u(x^2 + y^2))}{1 - 4u^2}\right)}{\sqrt{1 - 4u^2}} \quad \binom{m}{k} B_{n+k}$$

$$\frac{e^{-px} T_n(x)}{\sqrt{1-x^2}} \quad {}_2F_1\left(\begin{matrix} a & b \\ a+b+1/2 \end{matrix} \middle| z\right) \quad 2^n$$

$$\{T_n(x)\} \quad \text{“=”} \quad \begin{cases} (1-x^2) T_n''(x) - x T_n'(x) + n^2 T_n(x) = 0 \\ n T_{n+1}(x) + 2(1-x^2) T_n'(x) - n T_{n-1}(x) = 0 \\ (+ \text{ initial values}) \end{cases}$$

# Gessel Walks



Allowed steps:  $\rightarrow, \leftarrow, \nearrow, \searrow$

$$Q(x, y, t) = \sum_{n, i, j} q_{i, j, n} x^i y^j t^n$$

**Theorem.** (“Gessel’s conjecture”)

[Kauers-Koutschan-Zeilberger 2008]

$$Q(0, 0, t) = {}_3F_2\left(\begin{matrix} 5/6, 1/2, 1 \\ 5/3, 2 \end{matrix} \middle| 16 t^2\right).$$

**Theorem.** [Bostan-Kauers 2010]

$Q(x, y, t)$  is algebraic.

(Estimated size of min. poly: 30 Gb.)

**Needs fast algorithms!**

# Guessing

Given the **first terms** of a series  $y(z) = y_0 + y_1 z + y_2 z^2 + \dots$ ,

find **a differential operator  $L$**  of order  $< r$ ,

with polynomial coefficients of degree  $< d$ ,

such that  $L \cdot y(z) = O(z^\sigma)$ .

→ That's just linear algebra

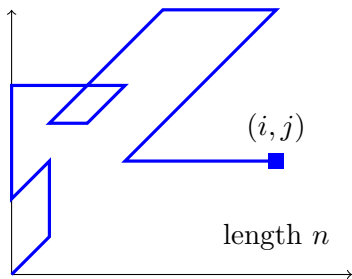
## Guess-and-prove

Faced with a conjectured identity  $A(z) = B(z)$ :

1. **compute** the initial terms of the series expansions of  $A$  and  $B$ ;
2. **guess** a common differential equation;
3. **verify** that it is satisfied and check a finite number of initial values.



# Gessel Walks



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(Estimated size of min. poly: 30 Gb.)

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# **3** **Approximations**

# Differential Equations and Recurrences

The **Taylor coefficients** of a D-finite function  $y(z) = \sum_{n=0}^{\infty} y_n z^n$  obey a linear **recurrence relation** with polynomial coefficients:

$$b_s(n) y_{n+s} + \cdots + b_1(n) y_{n+1} + b_0(n) y_n = 0.$$

(And conversely, for D-finite formal power series.)

**Example.**  $y'' + y = 0 \Rightarrow (n+1)(n+1)y_{n+2} + y_n = 0$

**More recurrences:** Generalized Fourier series.  $y(x) = \sum_{n \in \mathbb{Z}} c_n T_n(x)$

$$c_{n-3} + (2n-1)(2n-5)c_{n-1} - (2n+1)(2n+5)c_{n+1} - c_{n+3} = 0$$

# Recurrences Lead to Fast Algorithms

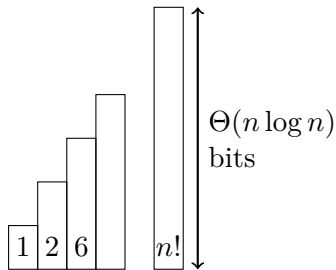
**First  $n$  terms**  $O(n)$  ring operations

**$n$ th term**  $\tilde{O}(\sqrt{n})$  ring operations (“baby steps-giant steps”)

**$n$ th term, over  $\mathbb{Q}$**   $O(M(n \log^2 n))$  bit operations (“binary splitting”)

## Naïve algorithm

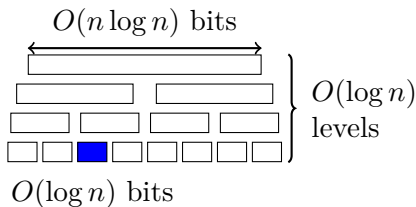
$$n! = n \cdot (n-1)!$$



**Time:**  $\Omega(n^2 \log n)$

## Binary splitting

$$n! = (n \cdots (\lfloor n/2 \rfloor + 1)) \cdot \lfloor n/2 \rfloor!$$



**Time:**  $O(M(n \log n) \log n)$

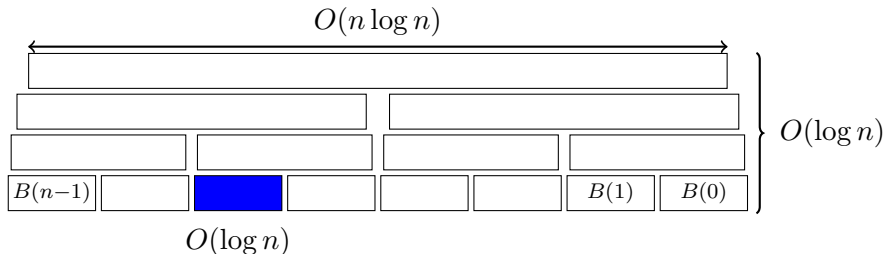
# Binary Splitting for D-Finite Series

Same idea, using a matrix product tree:

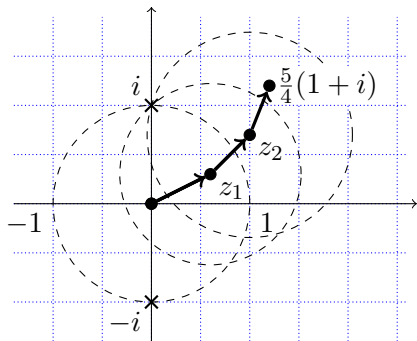
$$y(z) = \sum_{n=0}^{\infty} y_n z^n$$

$$\mathbf{S}_n = \sum_{k=0}^{n-1} y_k \zeta^k$$

$$\begin{bmatrix} y_{n+1} \\ \vdots \\ y_{n+s} \\ \mathbf{S}_{n+1} \end{bmatrix} = \begin{bmatrix} \color{cyan} B(n) \end{bmatrix} \begin{bmatrix} y_n \\ \vdots \\ y_{n+s-1} \\ \mathbf{S}_n \end{bmatrix}$$



# Numerical Analytic Continuation



$$\arctan\left(\frac{5}{4}(1+i)\right) ?$$

$$\begin{bmatrix} y(z_1) \\ y'(z_1) \end{bmatrix} = \begin{bmatrix} 1 & 0.57\dots + 0.22\dots \\ 0 & 0.72\dots - 0.20\dots \end{bmatrix} \begin{bmatrix} y(0) \\ y'(0) \end{bmatrix}$$

$$\begin{bmatrix} y(z_2) \\ y'(z_2) \end{bmatrix} = \begin{bmatrix} 1 & 0.36\dots + 0.32\dots \\ 0 & 0.75\dots - 0.07\dots \end{bmatrix} \begin{bmatrix} y(0) \\ y'(0) \end{bmatrix}$$

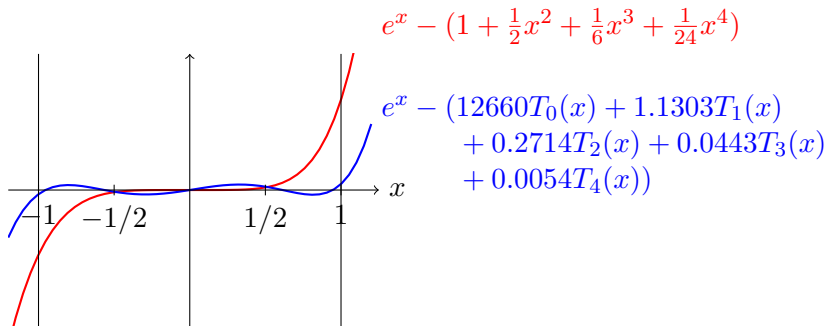
...

For a fixed function  $y$  and a fixed evaluation point  $z \in \mathbb{C}$ ,

one may compute  $y(z)$  with absolute error  $\leq 2^{-n}$

in time  $O(M(n \log^2 n))$  and space  $O(n)$ . [Chudnovsky<sup>2</sup>, v.d. Hoeven, M.]

# Polynomial Approximation on Intervals



$$y(x) = \sum_{n \in \mathbb{Z}} c_n T_n(x)$$

$$c_{n+1} + 2n c_n - c_{n-1} = 0$$

## Obstacles:

- high order  $\leftrightarrow$  divergent solutions
- ini. cond.  $\notin \mathbb{Q}$

## Solution:

[Benoit-Joldeş-M.]

- backward recursion (“Clenshaw’s algo”)

# Error Bounds

## Majorizing series (a priori bounds)

- Bound the differential equation with a simple “model equation”:

$$y'(z) = a(z) y(z) \quad \Leftarrow \quad g'(z) = \frac{1}{(1-\alpha z)} g(z)$$

- Solve the model equation and study the solutions:

$$\left| \sum_{k=n}^{+\infty} y_k z^k \right| \leq \sum_{k=n}^{+\infty} g_k |z|^k \leq ?$$

## Fixed-point bounds (a posteriori bounds)

- Rewrite as  $T(y) = y$ , say with  $\|T(f) - T(g)\| \leq \lambda \|f - g\|$ ,  $\lambda < 1$
- Given an approximate solution  $\tilde{y}$ , check that  $T(B(y, \varepsilon)) \subseteq B(y, \varepsilon)$



# Double-Precision Implementation

## The Problem

- Writing floating-point implementations of mathematical functions is hard!
- Special function libraries **performance** and **rigor** decades behind state-of-the-art elementary function libraries
- Libraries rarely provide “the right” function
  - $\exp(\sin(x))$
  - not accurate enough, more accurate than needed...

# Code Generation

$$x^2 Y_1'' + x Y_1' + (x^2 - 1) Y_1 = 0$$
$$Y_1(x) \sim -\frac{2}{\pi x} + \frac{x \ln x}{\pi} + \dots$$

as  $x \rightarrow 0$



```
double BesselY1 (double x)
{
    // generated code
}
```



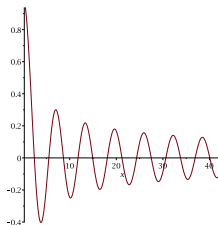
# A Prototype Code Generator

[Lauter-M. 2014]

**Input:** differential equation + initial values defining  $f$   
domain  $[a, b]$   
target accuracy  
(+ implementation constraints)

**Output:** C code

$$\left| \frac{\text{implem}(x) - f(x)}{f(x)} \right| \leq \varepsilon \text{ for all } x \in [a, b] \cap \text{double}$$



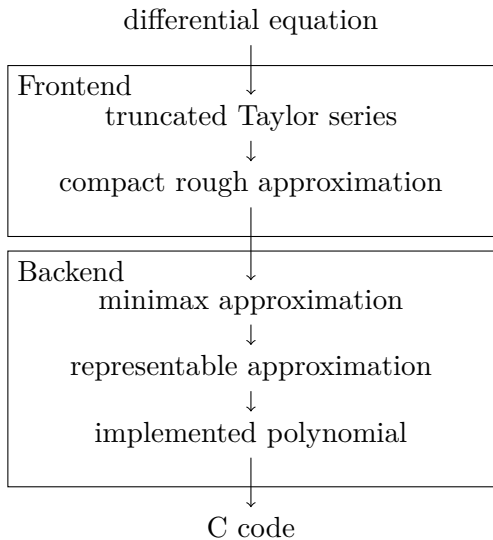
**Example:**

$$f(x) = J_0(x) \text{ (Bessel)}$$

$$[a, b] = [0.5, 42]$$

$$\varepsilon = 2^{-45}$$

# Architecture



- "Rigorous **polynomial approximation** pipeline"
- **Interesting points:**  
 $f(x) = 0$  or  $f(x) \approx 0$   
 $x = 0, x \rightarrow \infty$   
later: singularities
- Reduce to implementing well-chosen polynomials